Stochastically Constrained Best Arm Identification with Thompson Sampling

Le Yang cooperated with Siyang Gao, Cheng Li, Yi Wang

Email: le_yang@nus.edu.sg

Date: September 16, 2024

 Ω

4 ロ > 4 何 > 4

Contents

[Formulation](#page-7-0)

- [BFAI-TS Algorithm](#page-11-0)
- [Theoretical Results](#page-15-0)
- [Numerical Results](#page-20-0)

[Conclusions](#page-22-0)

KON YANK YENYEN EL MON

[Introduction](#page-2-0)

KOX KOX KEX KEX E VOQO

[Introduction](#page-2-0)

Best Arm Identification (BAI) Problem

 $x = x$

a affiliar

÷

書き

 Ω

Thompson Sampling (TS) Algorithm

TTTS is an asymptotically optimal algorithm for BAI problems [Russo \[2020\]](#page-24-0).

TS was originally proposed by Thompson in 1933 [\[Thompson, 1933](#page-24-1)] for the multi-armed bandit (MAB) problem.

TS has been extended to tackle a wide range of variant MAB problems. Combinatorial bandits [\[Sankararaman and Slivkins, 2018](#page-24-2)] Contextual bandits [\[Agrawal and Goyal, 2013](#page-24-3)] Online problems [\[Gopalan, Mannor, and Mansour, 2014\]](#page-24-4)

 Ω

イロト イ押ト イヨト イヨト

Best Feasible Arm Identification (BFAI) Problem

Le_yang@nus.edu.sg September 16, 2024 6/26

 \Rightarrow

∍

 QQ

Related Literature

重

 299

 $(1 - 1)$ $(1 - 1)$ $(1 - 1)$ $(1 - 1)$ $(1 - 1)$ $(1 - 1)$ $(1 - 1)$ $(1 - 1)$ $(1 - 1)$

[Formulation](#page-7-0)

KOX KOX KEX KEX E VOQO

Best Feasible Arm Identification (BFAI) Problem

- Goal: Identify the best arm: $I^* = \{i: argmax_i \mu_{i0} \ s.t. \mu_{ij} \leq \gamma_i, \forall j = 1, 2, \dots, m\}.$ \bullet
- Assumption: The best feasible arm is unique. Let $I^* = 1$. \bullet

 $\mathbf{A} \cap \mathbf{D} \times \mathbf{A} \cap \mathbf{B} \times \mathbf{A} \subset \mathbf{B} \times \mathbf{A} \times \mathbf{B} \times \mathbf{B} \times \mathbf{B}$

 QQ

Bayesian Framework

Normal prior distributions:

Posterior distribution:

$$
\begin{aligned} \Pi_t =& \mathcal{N}(\mu_{t,10}, \sigma_{t,10}^2) \otimes \ldots \otimes \mathcal{N}(\mu_{t,1m}, \sigma_{t,1m}^2) \otimes \ldots \\ \otimes & \mathcal{N}(\mu_{t,k0}, \sigma_{t,k0}^2) \otimes \ldots \otimes \mathcal{N}(\mu_{t,km}, \sigma_{t,km}^2). \end{aligned}
$$

 2990

メロメ メ御 トメ ミメ メ ミメー

Asymptotic Optimal

 \sim

The Contract of the Contract

EL 15 ÷

[BFAI-TS Algorithm](#page-11-0)

KORKAN KERKER DRAM

Process of Extending

Decision Policy

Algorithm Description

Algorithm Analysis

$$
\text{Define }\phi_{t,i}\triangleq\mathbb{P}(I_t=i|\mathcal{F}_{t-1})\text{ and }\bar{\phi}_{t,i}\triangleq\frac{\sum_{l=2}^{t}\phi_{l,i}}{t}.
$$

The probability of pulling arm i in round t :

$$
\phi_{t,i} = \frac{c_t}{k} + (1 - \beta)P_{t,i} \sum_{i' \neq i} \left(\frac{P_{t,i'}}{1 - P_{t,i'}} (1 - c_t) + \frac{c_t}{k - 1} \right) + P_{t,i} \beta (1 - c_t),
$$

where c_t is the probability that samples from all the arms are infeasible in round t .

$$
P_{t,1}\rightarrow 1, \, \phi_{t,1}\rightarrow \beta, \, c_t\rightarrow 0
$$

$$
\frac{P_{t,i}}{1-P_{t,1}} \to \frac{\phi_{t,i}}{1-\phi_{t,1}}, \ i \neq 1.
$$

∍

 Ω

K ロト K 御 ト K 君 ト K 君 ト

[Theoretical Results](#page-15-0)

 $x = x + y = x + z = 1$

 $= \Omega Q$

Notations

- ${\mathcal F}$ the set of feasible arms, i.e., ${\mathcal F} \triangleq \{i: \mu_{ij} \leq \gamma_j, \forall j \in \{1,2,...,m\}\};$
- I^* the best feasible arm, i.e., $I^*\triangleq \mathop{\mathrm{argmax}}_{i\in\mathcal{F}}\mu_{i0};$
- \mathcal{F}_w the set of feasible but suboptimal arms, i.e., $\mathcal{F}_w \triangleq \{i : i \in \mathcal{F} \text{ and } i \neq I^*\};$
	- \mathcal{I}_{b} the set of infeasible arms with objective performance no worse than I^{*} , i.e., ${\mathcal I}_b \triangleq \{i: \mu_{I^*0} \le \mu_{i0}$ and $\exists j \in \{1,2,\ldots,m\}$ such that $\mu_{ij} > \gamma_j\};$
- \mathcal{I}_w the set of infeasible arms with objective performance worse than I^* , i.e., ${\mathcal I}_w \triangleq \{i: \mu_{I^*0} > \mu_{i0}$ and $\exists j \in \{1,2,\ldots,m\}$ such that $\mu_{ij} > \gamma_j\};$
- \mathcal{M}_F^i the set of constraints estimated as satisfied by arm i in round t , i.e. ${\mathcal M}^i_F \triangleq \{j: \mu_{ij} \leq \gamma_j \text{ for } j \in \{1,2,\ldots,m\}\};$
- \mathcal{M}^i_I the set of constraints estimated as violated by arm i in round t , i.e. $\mathcal{M}_{I}^{i} \triangleq \{j: \mu_{ij} > \gamma_{j} \text{ for } j \in \{1,2,\ldots,m\}\}.$

∍

 Ω

イロト イ押 トイヨ トイヨト

Definitions

 $(\alpha_2^\beta,...,\alpha_k^\beta)$:

the optimal sampling rates of the remaining $k - 1$ arms, which satisfy the following optimality condition

$$
\sum_{i=2}^k \alpha_i^\beta = 1-\beta, \text{ and } \mathcal{R}_i = \mathcal{R}_{i'} \text{ for any } i\neq i' \neq 1,
$$

$$
\mathcal{R}_i = \tfrac{(\mu_{i0} - \mu_{10})^2}{(\sigma_{i0}^2/\alpha_i^{\beta} + \sigma_{10}^2/\beta)}\mathbf{1}\{i \in \mathcal{F}_w \cup \mathcal{I}_w\} + \alpha_i^\beta \sum_{j \in \mathcal{M}_I^i} \tfrac{(\mu_{ij} - \gamma_j)^2}{\sigma_{ij}^2}\mathbf{1}\{i \in \mathcal{I}_b \cup \mathcal{I}_w\}.
$$

 Γ_{β} : the optimal sampling rates given $\beta \in (0,1)$, where

$$
\Gamma_{\beta} = \min_{i \neq 1} \left(\frac{(\mu_{i0} - \mu_{10})^2}{2(\sigma_{i0}^2/\alpha_i^{\beta} + \sigma_{10}^2/\beta)} \mathbf{1}\{i \in \mathcal{F}_w \cup \mathcal{I}_w\} + \alpha_i^{\beta} \sum_{j \in \mathcal{M}_I^i} \frac{(\mu_{ij} - \gamma_j)^2}{2\sigma_{ij}^2} \mathbf{1}\{i \in \mathcal{I}_b \cup \mathcal{I}_w\}, \min_{j \in \mathcal{M}_F^1} \beta \frac{(\mu_{1j} - \gamma_j)^2}{2\sigma_{1j}^2} \right).
$$

ă.

 QQ

メロメメ 倒 メメ きょく ミメー

Theorem

For the BFAI-TS Algorithm, $\mathbb{E}[N_{\beta}^{\epsilon}]<\infty$ for any $\epsilon>0$, where

$$
N^\epsilon_\beta \triangleq \inf\{t\in\mathbb{N}:|\mu_{n,ij}-\mu_{ij}|\leq \epsilon \text{ and } |N_{n,i}/n-\alpha^\beta_i|\leq \epsilon, \forall i\in A \text{ and } n\geq t\}
$$

given $\beta \in (0,1)$. The sample allocations of the algorithm is asymptotically optimal in the sense that

$$
\lim_{n \to \infty} \frac{N_{n,i}}{n} \xrightarrow{p} \alpha_i^{\beta} \quad \forall i \in A,
$$

KO FRANCIS A BIF KORA

Theorem

The following properties hold with probability 1:

1. For any $\beta \in (0,1)$, Γ_{β} shows the fastest rate of posterior convergence that any algorithm allocating β proportion of the total samples to the best feasible arm can possibly achieve

$$
\limsup_{n \to \infty} -\frac{1}{n} \log(1 - P_{n,1}) \le \Gamma_{\beta} \tag{1}
$$

and the BFAI-TS Algorithm achieves this rate with

$$
\lim_{n \to \infty} -\frac{1}{n} \log(1 - P_{n,1}) = \Gamma_{\beta}.
$$
 (2)

2. The term $\Gamma_{\beta*}$ shows the fastest rate of posterior convergence that any BFAI algorithm can possibly achieve

$$
\limsup_{n \to \infty} -\frac{1}{n} \log(1 - P_{n,1}) \le \Gamma_{\beta^*} \tag{3}
$$

and when the β of the BFAI-TS Algorithm is set to β^* , the algorithm achieves the optimal rate with

$$
\lim_{n \to \infty} -\frac{1}{n} \log(1 - P_{n,1}) = \Gamma_{\beta^*}.
$$
 (4)

[Numerical Results](#page-20-0)

 $x = x + y = x + z = 1$

 $=$ Ω

[Numerical Results](#page-20-0)

Fig. 1. PFS and their sampling rates on selected arms (Experiment 1)

[Conclusions](#page-22-0)

ALLAMATA ATA TAGO

Demonstrate the extensibility of the TS algorithm by proposing the BFAI-TS algorithm.

Solve a significant and common class of constrained optimization problem, the BFAI problem.

The BFAI-TS algorithm is asymptotically optimal and exhibits impressive numerical performance.

 QQ

KO K K @ K K B K K

References

- S. Agrawal and N. Goyal. Thompson sampling for contextual bandits with linear payoffs. In *International Conference on Machine Learning*, pages 127–135, 2013.
- A. Gopalan, S. Mannor, and Y. Mansour. Thompson sampling for complex online problems. In *International Conference on Machine Learning*, pages 100–108. PMLR, 2014.
- D. Russo. Simple Bayesian algorithms for best arm identification. *Operations Research*, 68(6):1625–1647, 2020.
- K. A. Sankararaman and A. Slivkins. Combinatorial semi-bandits with knapsacks. In *International Conference on Artificial Intelligence and Statistics*, pages 1760–1770. PMLR, 2018.
- W. R. Thompson. On the likelihood that one unknown probability exceeds another in view of the evidence of two samples. *Biometrika*, 25(3-4):285–294, 1933.

 Ω

 $\left\{ \begin{array}{ccc} 1 & 0 & 0 \\ 0 & 1 & 0 \end{array} \right\}$, $\left\{ \begin{array}{ccc} 0 & 0 & 0 \\ 0 & 0 & 0 \end{array} \right\}$

Thank you!

START IN

a mille

 λ . The set

The Co \equiv

 \mathcal{A}

 Ω